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Extremal Configurations of Hinge Structures

Ciprian Borcea and Ileana Streinu

Abstract

We study body-and-hinge and panel-and-hinge chains in R^d , with two marked points: one on the first body, the other on the last. For a general chain, the squared distance between the marked points gives a Morse-Bott function on a torus configuration space. Maximal configurations, when the distance between the two marked points reaches a global maximum, have particularly simple geometrical characterizations. The three-dimensional case is relevant for applications to robotics and molecular structures.

Keywords: extremal configuration, Morse-Bott function, Hessian matrix, hinge structure, maximum reach, revolute-jointed manipulator.

AMS Subject Classification: 53A17

Introduction

This work is an extension of our study of singularities of hinge structures [BS]. We refer to that paper for basic notions and background.

The hinge structures considered here will be body-and-hinge or panel-and-hinge chains in R^d which have a point marked on the first body and a point marked on the last body. In dimension three, these hinge structures would model serial manipulators with revolute joints with a marked base-point and the end-effector as the other marked point¹. Likewise, the panel-and-hinge case may serve as model for “backbone” protein chains [BT, CP, BS]. Extremal configurations will be those where the squared distance function between the two marked points (origin and terminus, or “head” and “tail”) reaches a local maximum or minimum.

¹The geometrical models have no rotational limitations around the joints and no self-collision prohibitions.

Robotics is obviously concerned with extremal reaches of manipulators and the closely related problem of identifying the total workspace of a robot. A *necessary* condition for extremal configurations was recognized and proven in several papers [D, KW, SD, S]. In the words of [S], where the base-point may be chosen arbitrarily, and the end-point is called “hand”, this necessary condition says: “the line of sight from that point to the hand must intersect all turning axes”². However, all critical points with non-zero value for the squared distance function satisfy the condition, and they grow exponentially with the number of hinges.

In this paper, we refine the study of extremal configurations and obtain, in particular, a very simple *necessary and sufficient* characterization of the global maximum. It may be observed here that the approach used to identify the global maximum configuration (which is unique for generic body-and-hinge chains) cannot be fully adapted for the global minimum, although the panel-and-hinge case offers a fair degree of similarity (Theorems 10 and 11). The distinction, we suggest, stems from the possibility to reinterpret the global maximum as a global minimum of a related problem. Once recognized, the criterion for the global maximum can be proven with completely elementary means.

Global Maximum Theorem: *A body-and-hinge chain is in a global maximum configuration if and only if the segment from the origin to the end-point intersects all hinges in their natural order.*

The panel-and-hinge case has sufficient specificity to warrant separate treatment, most particularly in dimension three. Local extrema must be global extrema, but are not unique if not flat.

In the final section we discuss some variations.

1 A Morse-Bott function for chains with two marked points

In this section we consider body-and-hinge chains with $(n + 1)$ bodies and two marked points, one on the first and the other on the last body. The ambient dimension will be d , and we identify the first body with the fixed reference coordinate system $R^d \equiv B_1$, with the marked point at the origin. The point on the last body will be the end-point.

²This incidence of the origin-to-terminus line with the hinges is understood projectively, that is, includes the possibility of parallelism.

The composition of the *end-point map* $e : (S^1)^n \rightarrow R^d$ with the *squared norm function* $R^d \rightarrow R$ gives the *squared distance function* of the end-point to the origin:

$$F : (S^1)^n \rightarrow R, \quad F(\theta) = \langle e(\theta), e(\theta) \rangle \quad (1)$$

We'll use $T_n = (S^1)^n$ as another notation for the n -torus parametrizing the configuration space of our body-and-hinge chain.

The critical points of F are described by:

Proposition 1. *Let $n \geq d$. If zero is a value of F , then all points in $X_0 = F^{-1}(0) \subset T_n$ are critical points of F . The critical points with non-zero critical values are those configurations which have all hinges (projectively) incident with the line connecting the origin to the end-point.*

In the generic case, $F : T_n = (S^1)^n \rightarrow R$ is a Morse-Bott function, which has only isolated critical points for non-zero critical values, while the fiber over zero $X_0 \subset T_n$, when non-empty, will be smooth, of dimension $n - d$.

Proof: The squared norm on R^d has a critical point at the origin, hence all configurations with the two marked points coinciding (i.e. figuratively, when “the head bites the tail”) will be critical for F .

For non-zero critical values, the argument is similar to the one used in [BS] : at a critical configuration, rotating the part of the chain from hinge A_i on, as a rigid piece, must preserve, infinitesimally, the (squared) distance “head-to-tail”, that is: must produce a velocity vector for the end-point orthogonal to the line between the marked points. That requires the (projective) incidence of hinge and line.

In the generic case, the origin will be a regular value of the end-point map, and, when zero is a value of F , it will give a smooth (not necessarily connected), codimension d fiber $X_0 = F^{-1}(0) \subset T_n$. The fact that the remaining critical points are isolated will follow from the examination of the corresponding Hessian. The Bott non-degeneracy condition [G] along (the connected components of) X_0 will be verified at that stage as well.

1.1 Set-up for computing the Hessian matrix

In a given configuration, hinge number i will be determined by the vector t_i which is the perpendicular projection of the origin on the corresponding hinge,

plus the normal direction to the hyperplane formed by the hinge and the origin. One may keep track of global orientations, but for our computations, a local choice of unit normal ν_i will suffice. Thus $(t_i, \nu_i) = (t_i(\theta), \nu_i(\theta))$ ‘encodes’ the i^{th} hinge.

At any critical point, we may assume the $\theta \in (S^1)^n$ labelling of the configuration space introduced by the following rule: the critical point is $\theta = 0$, and the position for arbitrary $\theta = (\theta_1, \dots, \theta_n)$ is obtained by rotating the last body around the last hinge with angle θ_n , then rotating the last two bodies (as a rigid piece) around the last but one hinge with angle θ_{n-1} , and so on until, at last, the whole (rigid) piece thus formed with all the bodies from the second to the last is rotated around the first hinge with angle θ_1 . All rotations, for $i = n, \dots, 1$, are using the sense dictated by a fixed orientation, say $\{\nu_i, t_i\}$ in the vector plane $[\nu_i, t_i]$ they span.

Let $R_i(\omega)$ stand for the *linear operator in R^d* which gives the *rotation* with angle ω around $[\nu_i, t_i]^\perp$, with the orientation fixed as above. Then, $\frac{\partial R_i}{\partial \omega}(\omega)$ is a skew-symmetric operator vanishing on $[\nu_i, t_i]^\perp$, and we put $\frac{\partial R_i}{\partial \omega}(0) = S_i$.

With a second derivation $\frac{\partial^2 R_i}{\partial \omega^2}(\omega) = -P_i R_i(\omega)$ with P_i denoting the orthogonal projection on the 2-subspace $[\nu_i, t_i]$. Thus: $\frac{\partial^2 R_i}{\partial \omega^2}(0) = -P_i$

We let $x \in R^d$ denote the position of the end-point for the critical configuration under investigation. With the parametrization and notation just described, and the abbreviation $R_i(\theta_i) = R_i$, the *end-point function* is:

$$e(\theta) = R_1 R_2 \dots R_n x + R_1 \dots R_{n-2} (I - R_n) t_n + \dots + R_1 (I - R_2) t_2 + (I - R_1) t_1 \quad (2)$$

This gives:

$$e(0) = x, \quad \frac{\partial e}{\partial \theta_i}(0) = S_i(x - t_i)$$

$$\frac{\partial^2 e}{\partial \theta_i^2}(0) = P_i(t_i - x), \quad \frac{\partial^2 e}{\partial \theta_i \partial \theta_j}(0) = S_j S_i(x - t_i), \quad j < i$$

Considering that $S_i t_i = -||t_i||\nu_i$, the resulting entries for the Hessian matrix are:

$$\frac{1}{2} \frac{\partial^2 F}{\partial \theta_i \partial \theta_j}(0) = \langle S_i(x - t_i), S_j(x - t_j) \rangle + \langle S_j S_i(x - t_i), x \rangle =$$

$$\begin{aligned}
& = \langle S_i(x - t_i), S_j(x - t_j) \rangle - \langle S_i(x - t_i), S_j x \rangle = \langle S_i(t_i - x), S_j t_j \rangle = \\
& = \langle S_i(1 - \frac{\langle x, t_i \rangle}{\langle t_i, t_i \rangle}) t_i, S_j t_j \rangle = (1 - \frac{\langle x, t_i \rangle}{\langle t_i, t_i \rangle}) \langle S_i t_i, S_j t_j \rangle = \\
& = (1 - \frac{\langle x, t_i \rangle}{\langle t_i, t_i \rangle}) \|t_i\| \cdot \|t_j\| \langle \nu_i, \nu_j \rangle
\end{aligned}$$

for $j \leq i$.

We retain the result of this computation as:

Proposition 2. *With adequate parametrization, the symmetric $n \times n$ Hessian matrix for the squared end-point distance function F at a critical configuration $\theta = 0$, with end-point at $e(0) = x$, has entries:*

$$\frac{1}{2} \frac{\partial^2 F}{\partial \theta_i \partial \theta_j}(0) = (1 - \frac{\langle x, t_i \rangle}{\langle t_i, t_i \rangle}) \|t_i\| \cdot \|t_j\| \langle \nu_i, \nu_j \rangle, \quad j \leq i \quad (3)$$

and, after the change of basis $e_i \mapsto \frac{1}{\|t_i\|} e_i$, corresponds with the quadratic form given by:

$$h_{ij} = h_{ji} = (1 - \frac{\langle x, t_i \rangle}{\langle t_i, t_i \rangle}) \langle \nu_i, \nu_j \rangle, \quad j \leq i \quad (4)$$

For a non-zero critical value, the coefficients $\alpha_i = \frac{\langle x, t_i \rangle}{\langle t_i, t_i \rangle}$ are obtained geometrically from the intersections of the line through the origin and the end-point x with the hinges, since these intersections are precisely the points $\frac{1}{\alpha_i} x$, $i = 1, \dots, n$.

When F takes the value zero, the Hessian at a critical point in $X_0 = F^{-1}(0)$ is equivalent to the Gram matrix $h_{ij} = \langle \nu_i, \nu_j \rangle$ of the normals, which is semi-positive definite and has, in the generic case, rank d . Thus, the null-space corresponds to the tangent space of X_0 at the critical point under consideration. This is the Bott non-degeneracy property required along the critical manifold X_0 .

1.2 An upper-bound for the number of isolated critical points

An upper-bound for the number of isolated critical points can be obtained from the following *complexification*: $T_n = (S^1)^n$ is complexified to $(P_1(C))^n$ by considering each circle S^1 as the real locus of the corresponding complex conic:

$$P_1(C) \approx \{x \in P_2(C) : x_1^2 + x_2^2 = x_0^2\}$$

$$\frac{x_1}{x_0} = \cos \theta_i, \quad \frac{x_2}{x_0} = \sin \theta_i$$

With some choice of a reference simplex in each hinge $A_i(\theta)$, say $a_1^i(\theta), \dots, a_{d-1}^i(\theta)$, the condition that the ‘head-to-tail’ line meets this hinge becomes

$$\det[a_1^i(\theta) \dots a_{d-1}^i(\theta) e(\theta)] = 0$$

This defines in $(P_1)^n$ a hypersurface of multi-degree $2(d, \dots, d, 1, \dots, 1)$, with positions d up to the i^{th} coordinate.

The intersection of these n hypersurfaces yields, in the complex count, the number:

$$2^n \int_{(P_1)^n} (h_1 + \dots + h_n)(dh_1 + h_2 + \dots + h_n) \dots (dh_1 + \dots + dh_{n-1} + h_n) = 2^n \sum_{k=0}^{n-1} A(n, k) d^k$$

where h_i stands for the class of a point in the i^{th} factor P_1 , and $A(n, k)$ denote Eulerian numbers. Thus, we have:

Proposition 3. *The number of isolated critical points for $F : T_n \rightarrow R$ is bounded by*

$$2^n \sum_{k=0}^{n-1} A(n, k) d^k \tag{5}$$

1.3 A Meyer-Vietoris sequence (or Morse-Bott theory)

The determination of critical configurations and their indices can be used in the following setting: we *assume* a generic body-and-hinge chain, $n \geq 3$ and zero to be a value of F . This ensures the smoothness of the $(n-3)$ -dimensional fiber $X_0 = F^{-1}(0)$ over zero, and, for small enough $\epsilon > 0$, a diffeomorphism $X_0 \times B_\epsilon^3 \approx F^{-1}[0, \epsilon]$, where B_ϵ^3 stands for the 3-dimensional ball of radius ϵ . The $(n-1)$ -dimensional fiber $X_\epsilon = F^{-1}(\epsilon)$ is thereby identified with $X_0 \times S_\epsilon^2$.

We'll use the notation $T_n = (S^1)^n$ for the n -dimensional torus representing the configuration space of our chain with $(n+1)$ panels and n hinges. For small and nearby values $0 < \gamma < \epsilon < \delta$ we put:

$$U = F^{-1}[0, \delta), \quad V = T_n \setminus F^{-1}[0, \gamma]$$

assuming all non-zero critical values of F greater than δ . Thus, the two open sets cover the torus T_n , and we have *homotopy equivalences*:

$$U \sim X_0, \quad U \cap V \sim X_\epsilon = X_0 \times S^2$$

The corresponding Meyer-Vietoris exact sequence gives:

$$\dots \rightarrow H_i(X_0 \times S^2) \rightarrow H_i(X_0) \oplus H_i(V) \rightarrow H_i(T_n) \rightarrow H_{i-1}(X_0 \times S^2) \rightarrow \dots$$

and in particular, we have the relation of Euler-characteristics:

$$e(X_0 \times S^2) + e(T_n) = e(X_0) + e(V) \quad \text{i.e.} \quad e(X_0) = e(V)$$

By Morse theory [M], the last Euler-characteristic can be expressed in terms of critical points as follows: put c_i for the number of critical points of index i ; then:

$$e(V) = \sum_{i=0}^n (-1)^{n-i} c_i$$

Note: The common convention would be to use $-F$ as the Morse function. Since we go with F , the signs are as above.

We obtained:

Proposition 4. *Suppose $n \geq 3$ and let the origin be a regular value of the end-point map for an otherwise generic chain. Then, the $(n - 3)$ -dimensional manifold X_0 parametrizing all configurations with the origin coinciding with the end-point has the Euler number:*

$$e(X_0) = \sum_{i=0}^n (-1)^{n-i} c_i \quad (6)$$

where c_i is the number of critical points of index i for F on $T_n \setminus X_0$.

1.4 The index of the Hessian matrix: panel-and-hinge case

The restriction to panel-and-hinge chains brings new structural aspects. We note first the presence of natural transformations of the configurations space $T_n = (S^1)^n$:

Transformations: Given any configuration, one may consider the hyperplane through one marked point and some hinge, and reflect the part of the chain from point to hinge in this hyperplane. The resulting transformation is obviously its own inverse i.e. an *involution*. Note that we may always reposition the structure with its first panel in its fixed location. Since the composition of reflecting in the first and then last panel gives a global rotation of the chain, these two operations represent the same transformation of T_n and this gives $2n - 1$ involutions on the configuration space, all transforming the fibers of F to themselves. Two such involutions commute when they implicate the same marked point or the respective portions of the chain do not overlap.

We have seen that, in case F reaches 0, $F^{-1}(0)$ is part of the critical locus, but all critical points for non-zero critical values are isolated in the generic case. The following definition refers to these isolated critical configurations corresponding to non-zero critical values of F :

Definition 5. *The 2^n flattened configurations when all panels lie in the same hyperplane (i.e. the codimension one subspace of the first panel, which is identified here with $x_d = 0$ in R^d) will be called flat critical configurations (points), while critical configurations which do not have all panels in the same hyperplane will be called non-flat critical configurations (points).*

It will be observed that, in a non-flat critical configuration, a new panel normal direction requires two consecutive hinges to meet the line from the origin to the

end-point in the same point and we'll call such a point a fold point (on the "head-to-tail" line). Fold points will be considered as labelled by the corresponding pair of consecutive hinges and ordered via this labelling.

Note: Flat critical configurations are fixed points of all involutions described above and the group they generate, while non-flat critical configurations have an orbit of cardinality 2^c under this group, where c is the number of fold points on the "head-to-tail" line. Assuming a generic case, $c + 1$ will give the number of distinct hyperplanes determined by origin and hinges .

We settle first the case of all flat critical configurations. Then all normal directions are the same and (regardless of the choice $\pm n_i$) Proposition 2 gives the signature as that of the (quadratic form with) matrix:

$$h_{ij} = h_{ji} = 1 - \frac{\langle x, t_i \rangle}{\langle t_i, t_i \rangle} = \beta_i, \quad j \leq i \quad (7)$$

Lemma 6. *The signature of the above matrix is determined by the signs of:*

$$\beta_1 - \beta_2, \beta_2 - \beta_3, \dots, \beta_{n-1} - \beta_n, \beta_n$$

Proof: The symmetric matrix $H = (h_{ij})$ is the expression of the quadratic form $q(x) = x^t H x$ in the standard basis $e_i, i = 1, \dots, n$. If we do the (unimodular) change of basis:

$$e_i \mapsto \tilde{e}_i = e_i - e_{i+1}, \quad i = 1, \dots, n-1; \quad e_n = \tilde{e}_n$$

we obtain a diagonal matrix with the indicated entries.

As mentioned in Proposition 2, if we mark the end-point by x , and the intersections of the "head-to-tail" line with the k^{th} hinge is $a_k x$, then $\beta_k = 1 - a_k^{-1}$. For a generic chain in a flattened position, hinges will intersect the "head-to-tail" line in different points, and we'll have a non-degenerate Hessian. Since by definition, the index of a symmetric matrix is the number of negative eigen-values, we have:

Proposition 7. *The index of a flat critical configuration with hinges meeting the line from the origin to the end-point x at $a_k x$, $k = 1, \dots, n$, is the number of negative values in the list:*

$$1 - \frac{1}{a_n}, \frac{1}{a_n} - \frac{1}{a_{n-1}}, \dots, \frac{1}{a_2} - \frac{1}{a_1} \quad (8)$$

Corollary 8. *A flat critical configuration is a local maximum if and only if the hinges meet the oriented segment $(0, x)$ in points $a_k x$, $k = 1, \dots, n$, lined-up in their natural order, that is:*

$$0 < a_1 < a_2 < \dots < a_n < 1 \quad (9)$$

For a convenient geometric formulation of the existence of a local minimum in a flat configuration we conceive of the “head-to-tail” line as completed to a projective line, and the complement of the affine segment $[0, x]$ gives then the open arc from 0 to x passing through the “point at infinity”.

Corollary 9. *A flat critical configuration is a local minimum if and only if the hinges meet the oriented arc from 0 to x passing through the point at infinity in their natural order. This means one the following:*

$$a_n < \dots < a_2 < a_1 < 0 \quad \text{or} \quad (10)$$

$$a_k < \dots < a_1 < 0 < 1 < a_n < \dots < a_{k+1} \quad \text{or} \quad (11)$$

$$1 < a_n < \dots < a_2 < a_1 \quad (12)$$

Remark: In fact, the projective formulation allows some relaxation in the genericity assumptions and one of the hinges may be parallel to the “head-to-tail” line in that flat configuration.

1.5 Extremal configurations

In this section we refine our description of extremal configurations of panel-and-hinge chains in R^d with two marked points.

We’ll elaborate on our Corollaries 8 and 9 and address possible maxima and minima at non-flat critical configurations. One should remain aware of the involutions described in subsection 1.3. Recall that a fold point on the “head-to-tail” line is common to two consecutive hinges, say $a_k x = a_{k+1} x$. When speaking of the ordering of intersections of hinges with the “head-to-tail” line, either ordering may be envisaged for the two hinges, but we intend the ordering requested in the statement. The function F is the squared distance from “head” to “tail”.

Theorem 10. *A local maximal configuration for F is characterized by the fact that all hinges intersect the oriented segment $(0, x)$ in the natural order. Moreover, a local maximum is in fact the global maximum and is unique modulo the natural transformations generated by the involutions described above. Thus there are 2^μ maximal configurations equivalent under natural transformations, where μ is the number of fold points for any and all of them.*

Proof: Let us call *flat subsystem* in a non-flat critical configuration the panel-and-hinge structure obtained by retaining the consecutive axes which lie in the same hyperplane through the “head-to-tail” line i.e. the hinges corresponding to a specific normal direction ν_i (up to the first occurrence of a fold point, or between a fold point and the next, or up from the last fold point). The “head” and “tail” are inherited from the full configuration. Then the expression (4) we obtained for the Hessian shows that a local maximum requires all flat subsystems to be flat local maxima. Thus, by Corollary 8, all hinges intersect the oriented segment $(0, x)$ in the natural sequential order.

As a consequence, we may trace a “red line” on consecutive panels by following the segment $[0, x]$ in our local maximum configuration. Thus, any other configuration will display the “red line” as a polygonal arc from “head” to “tail” proving that our local maximum is the global maximum. It also follows that any other local maximum must have exactly the same pattern of planar subsystems and therefore be obtained from our maximal configuration by some composition of involutions.

Theorem 11. *A local minimal configuration for a non-zero value of F is characterized by the fact that all hinges intersect the oriented projective arc from 0 to x passing through the point at infinity (i.e. the complement of $[0, x]$) in the natural order. Moreover, if such a local minimum exists, it is in fact the global minimum and 0 is not a possible value for F . When 0 is not a possible value for F , all minima are achieved at 2^ν minimal configurations equivalent under natural transformations, where ν is the number of fold points for any and all of them.*

Proof: Simple adaptation of the “red line” argument presented above.

2 The global maximum as a global minimum

The criteria obtained above for extremal configurations of marked panel-and-hinge chains in R^d offer obvious suggestions for the more general *body-and-hinge case*.

It is an immediate observation that a marked body-and-hinge chain for which the segment from the origin to the end-point intersects all hinges in their natural order is in a global maximum configuration, since for any other configuration, the previous segment (drawn as a “red segment”) becomes a polygonal arc longer than the new segment from the origin to the end-point. What is less immediate is that any body-and-hinge chain actually reaches its global maximum in such a configuration. This will be proven by relating the global maximum to a global minimum.

Theorem 12. *Let a body-and-hinge chain be presented in a fixed configuration, with the origin as the marked point of the first panel, e the end-point (on the last body) and hinges given by codimension two affine subspaces A_i , $i = 1, \dots, n$. Consider variable points on each hinge $a_i \in A_i$ and the length of the polygonal arc going from the origin to the end-point through the points a_i in their natural order:*

$$f(a_1, \dots, a_n) = \|a_1\| + \|a_2 - a_1\| + \dots + \|a_n - a_{n-1}\| + \|e - a_n\| \quad (13)$$

The global maximum distance between the marked points of the given body-and-hinge chain equals the global minimum of the function f .

Proof: Note that f is continuous and a global minimum always exists.

The proof will follow from the simple case of a single hinge ($n=1$). In this case, a_1 must be the unique point of A_1 which allows a rotation of the segment from a_1 to e around A_1 to become a continuation of the segment from the origin to a_1 . Repeating this observation with respect to a_{i-1}, a_{i+1} and the hinge A_i , shows that the chain can be reconfigured so that the polygonal arc realizing the global minimum becomes a straight segment from the origin to the end-point intersecting all hinges in their natural order. That is the global maximum.

We have proven at the same time our:

Global Maximum Theorem: *A body-and-hinge chain is in a global maximum configuration if and only if the segment from the origin to the end-point intersects all hinges in their natural order. For a generic body-and-hinge chain, this global maximum is unique.*

Remarks: (i) In general, a body-and-hinge chain may well have local maxima which are not global maxima. By (4), the segment from the origin to the end-point must intersect all hinges.

(ii) f is continuous, but not differentiable when $a_i = a_{i+1}$ for some i .

Examples show that global minima for body-and-hinge chains do not have to respect the pattern observed for panel-and-hinge chains. We may consider, for instance, just two hinges (say in R^3) and take the second body (containing these hinges) as reference. Then the marked points trace circles, each around a hinge. Examples of global minima with the hinges intersecting in the order A_2, A_1 the projective arc from origin to the endpoint passing thorough infinity, can be easily produced. What still holds true, by (4), for any local minimum is the fact that all hinges intersect the projective arc from the origin to the end-point, passing through the point at infinity.

3 Variations on the same theme

A few ‘variations’ of these techniques should be mentioned before concluding. Again, the issues are important in robotics and the necessary conditions have been detected in the literature [SR, SD].

If we abandon the first marked point, we’ll rather be concerned with the squared distance from the end-point to the first hinge. The resulting critical configurations for non-zero values will require the *perpendicular* from the end-point to the first hinge to meet (projectively) all hinges. The global maximum will require the natural ordering of these intersections on the segment from the foot of the perpendicular to the end-point.

In dimension three, a similar scenario holds if we abandon both marked points and only inquire about the squared distance between the first hinge and the last hinge. Critical configurations for non-zero values will have the common perpendicular of these two hinges intersect (projectively) all the intermediate hinges. The global maximum will require the natural ordering.

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